

# The design and implementation of a benchmark based on an iterated post-selective protocol

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# The team, collaborators and companies providing qc time

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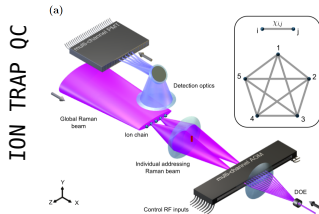
N. Linke



A. Green



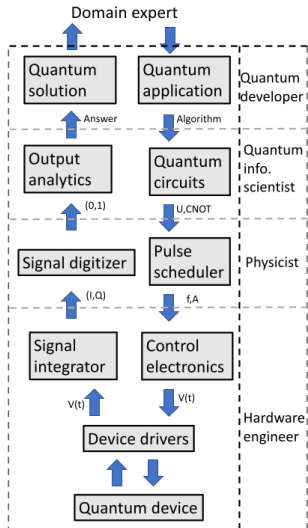
E. Mossman



IBM Q™



# Current cloud enabled quantum computing



←  
This talk is in this level

## Quantum benchmark

- ☞ Why a quantum benchmark? How to design a quantum benchmark?
- ☞ Implementation: Decompositions. Choosing the protocol. What to measure?
- ☞ Results



## Quantum benchmark: Why?

We have imperfect, highly noisy quantum computers (NISQ-era), so what to do with them? How to decide:

- Which technology is more promising (investment \$\$\$)?
- Is a single physical system (e.g. superconducting qubits, trapped ions, photonic chips) *the* universal system? Is it that a technology is suitable for a specific task?
- Which types of tasks are worth investigating (inside the current qc-capabilities)?
- (etc.)


## Quantum benchmark: Why?

An early type of benchmark/testing can help  
to decide these matters!

- Ok, but how to?
- Is a single number enough? In general no! Not even preferred in classical benchmarking.
- What are the characteristics of a good benchmark? [Maybe we can use what is already known in classical benchmarking...]

## Quantum benchmark: How to?

The current research divides roughly into two categories:



Hallmark protocols

(QC circuits or QI protocols,  
Quantum chemistry)

Randomized benchmarking

(QV, Mirror circuits)

Our testing/benchmarking fits in the first set.

## Quantum benchmark: How to?

Focusing in a *performance benchmark* (a benchmark in general also include security or energy-consumption metrics), for which performance of NISQ systems is usually investigated, it must fulfill the following characteristics<sup>1</sup>:

- The workload has to be representative of real-world applications.
- Exercise all critical services provided by the platforms (for a qc, both hardware and software).
- Not be tuned/optimized for a specific product [or likewise to tune a device for a specific task, i.e. QV and Honeywell <sup>2</sup>].
- Generate reproducible results.
- Not have any inherent scalability limitations [QV is not scalable!].

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<sup>1</sup> M. Vieira, H. Madeira, K. Sachs, and S. Kounev. Resilience benchmarking. page 283. Springer, 1 edition, 2012.

<sup>2</sup> <https://www.honeywell.com/us/en/news/2021/07/>

### The UnitaryForge package

- Quantum circuits are not written in general as a combination of one-qubit gates and CNOTs.
- A reliable, agnostic, decomposition is needed.
- The UnitaryForge package aims to the solution of this problem.
- Arbitrary, entangling two-qubit gate  $U (\in \text{SU}(4)) \rightarrow$  One-qubit gates and CNOTs.

Collab./Developer:

Y. Alonso

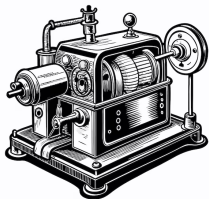


[github.com/Yeas3/UnitaryForge](https://github.com/Yeas3/UnitaryForge)

Deployment soon!

# Quantum benchmark: Decompositions

Get  $\vec{k} \in \mathbb{R}^3$



Khaneja-Glaser

$$X = (A_0 \otimes A_1) \exp \{-i\vec{k} \cdot \vec{\Sigma}\} (B_0 \otimes B_1),$$

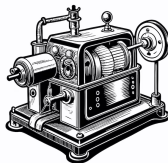
$$k_j \in \mathbb{R},$$

$$\Sigma_j = (\sigma_j \otimes \sigma_j) \quad \text{and}$$

$$A_i, B_i \in SU(2) \quad (i = 0, 1)$$

PRA 63, 032308 (2001)

Reorder, rewrite



Vidal-Dawson

$$\pi/4 \geq k'_x \geq k'_y \geq |k'_z| > 0 \rightarrow \times 3 \text{ CNOTs},$$

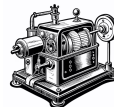
$$k'_x \geq k'_y > 0, k'_z = 0 \rightarrow \times 2 \text{ CNOTs},$$

$$\vec{k}' = (\pi/4, 0, 0)^T \rightarrow \times 1 \text{ CNOT}$$

+1q gates

PRA 69, 010301 (2004)

Factorize



Cayley

$$M^\dagger (\mathcal{C} \otimes \mathcal{D}) M \in SO(4) \rightarrow \mathcal{C}, \mathcal{D} \in SU(2)$$

IEEE Trans. Rob. 30, 1037 (2014)

```
U= [[ 0.12787+0.55181j -0.19527-0.59726j 0.29605+0.19217j 0.34233-0.20624j]
     [-0.12972+0.06977j -0.20892-0.2803j 0.34759-0.3652j -0.351 +0.69188j]
     [-0.54968+0.19556j -0.63534+0.22603j -0.15539+0.38409j -0.17914-0.03313j]
     [-0.29599-0.47902j -0.08102+0.147j 0.6661 +0.0766j 0]
```

Descomposición en cnots:

Número de cnots= 3

```
u2= [[-0.25669+0.65887j -0.25669+0.65887j]
     [ 0.25669+0.65887j -0.25669-0.65887j]]
```

```
v2= [[0.75155-0.65967j 0. +0.j ]
     [0. +0.j 0.65967-0.75155j]]
```

```
u3= [[0.-0.70711j 0.-0.70711j]
     [0.-0.70711j 0.+0.70711j]]
```

```
v3= [[0.98417+0.17721j 0. +0.j ]
     [0. +0.j 0.98417-0.17721j]]
```

```
w= [[0.70711+0.j 0. -0.70711j]
     [0. -0.70711j 0.70711+0.j ]]
```

Algoritmo KG:

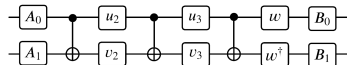
```
A0= [[ 0.41141+0.78943j -0.21083+0.40384j]
     [ 0.21083+0.40384j 0.41141+0.78943j]]
```

```
A1= [[-0.96144+0.1404j 0.17528+0.15877j]
     [-0.17528+0.15877j -0.96144-0.1404j ]]
```

```
B0= [[ 0.46875+0.12124j 0.74859-0.45297j]
     [-0.74859-0.45297j 0.46875-0.12124j]]
```

```
B1= [[-0.46875-0.26441j 0.54108-0.64621j]
     [-0.54108-0.64621j -0.46875+0.26441j]]
```

```
k= [-0.41389971 0.06501331 -0.17815323]
k0= 0.7853981634
```



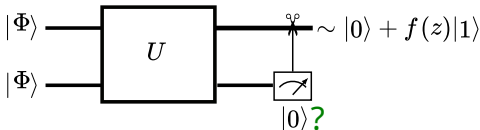
# Quantum benchmarks: The protocol

## Quantum state matching using a superattractive map

[O. Kálmán and T. Kiss, Phys. Rev. A 97, 032125 (2018)]

One step, 2 qubits. Initial state:

$$|\Phi\rangle \propto |0\rangle + z|1\rangle \rightarrow |\Psi_0\rangle = |\Phi\rangle \otimes |\Phi\rangle.$$



$U$  is an entangling unitary and  $f(z)$  is a quadratic rational complex function:

$$f(z) = \frac{a_0 + a_1 z + a_2 z^2}{b_0 + b_1 z + b_2 z^2}.$$

Collabs:

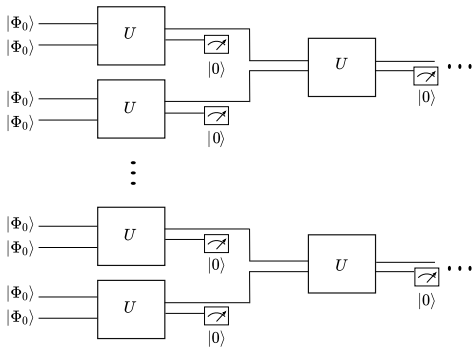
T. Kiss, O. Kálmán & Z. Udvarnoki



Full story: Phys. Scr. 98, 024006 (2023)

## Quantum benchmarks: The protocol

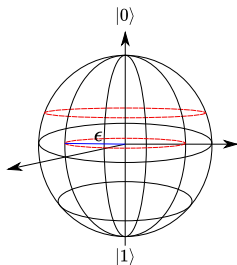
After  $n$ -steps:



$$\dots = \text{Unitary } U \text{ acting on } |0\rangle \text{ and } f^{(n)}(z)|1\rangle \sim |0\rangle + f^{(n)}(z)|1\rangle$$



# Quantum benchmarks: The protocol of quantum state matching

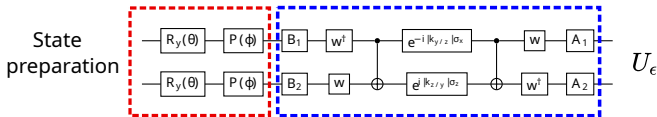


Transformation (not  $\phi$  dependant!)

$$f_{\epsilon}(z) = \frac{z^2}{\epsilon}, \quad z = e^{i\phi} \tan(\theta/2)$$

achieved using

$$U_{\epsilon} = \begin{pmatrix} \epsilon & -\frac{1}{\sqrt{2}}\sqrt{1-\epsilon^2} & \frac{1}{\sqrt{2}}\sqrt{1-\epsilon^2} & 0 \\ 0 & \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} & 0 \\ 0 & 0 & 0 & 1 \\ \sqrt{1-\epsilon^2} & \frac{1}{\sqrt{2}}\epsilon & -\frac{1}{\sqrt{2}}\epsilon & 0 \end{pmatrix}.$$



$w = (\mathbb{1} - i\sigma_x)/\sqrt{2}$ , see [G. Vidal and C. M. Dawson, Phys. Rev. A, 69:010301 (2004)].

## Characteristics:

- Super-attractive = Fast convergence.
- Orthogonalizing map for two states.
- $n$ -steps uses  $2^n$  qubits (good for testing purposes).
- **No classical overhead** (c.f. Quantum Volume benchmark).

- **Goal:** Benchmarking of quantum computers using the quantum state matching protocol.
- **Errors in the measurement:** statistical and device specific. A simple approach is to consider

$$\Delta_S/\sqrt{M} + \Delta_D,$$

Quantum benchmarks:  
What to measure?

$M$ : total number of experiments,  $\Delta_S$ : statistical noise,  $\Delta_D$ : device noise<sup>3</sup>.

Success probability of the post-selected state

$$\begin{aligned} p_s &= P_{|0\rangle} + P_{|1\rangle} \\ &= \epsilon^{2^{n+1}-2} \cos(\theta/2)^{2^{n+1}} + \sin(\theta/2)^{2^{n+1}}, \end{aligned}$$

at step  $n$  “filters out” the statistical error if we consider the interval

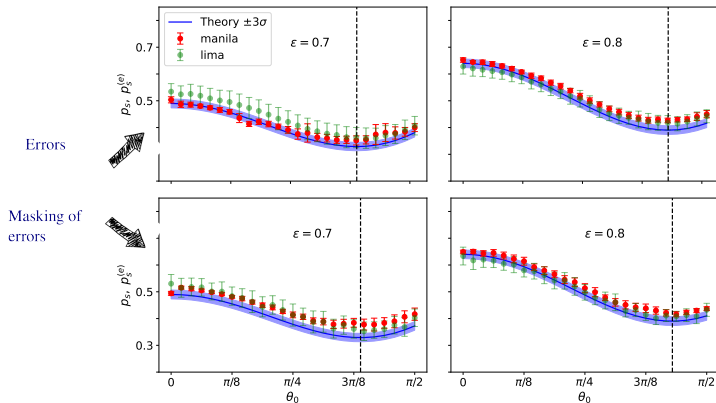
$$p_{s,\pm}^{(e)} = p_s \pm m\sigma,$$

for  $m$  an integer and  $\sigma = \sqrt{p_s(1-p_s)/M}$ .

<sup>3</sup>A. Cornelissen, et. al. arxiv.2104.10698

# Quantum Benchmark: Results I - Initial states

Influence of the initial (random) state for two different qcs (QVs )



**Figure:** One step, manila 32QV and lima 8QV. Upper row: deterministic  $\phi_0$  initial states. Bottom row: Random  $\phi_0$  initial states.

## Quantum benchmark: Results II - Coherent errors

The following set of errors are related with the implementation of gates in a certain qc, known as miss-rotations. The CNOT gate is the main source of such errors. In the IBM devices, the two-qubit interaction is the cross-resonance gate ( $\Lambda$  is the error in the rotation)

$$U(1/2 + \Lambda) = \exp(-i(1/2 + \Lambda)\pi Z_1 X_2/2)$$

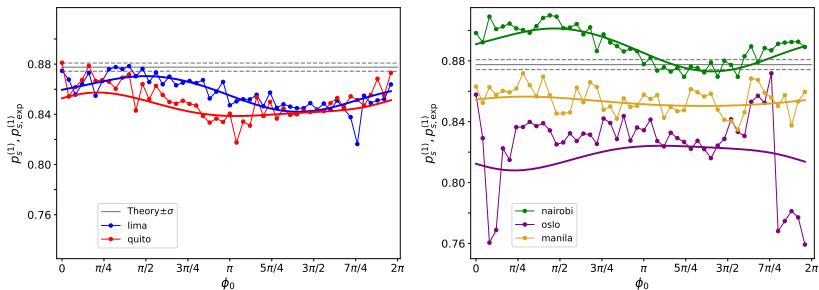
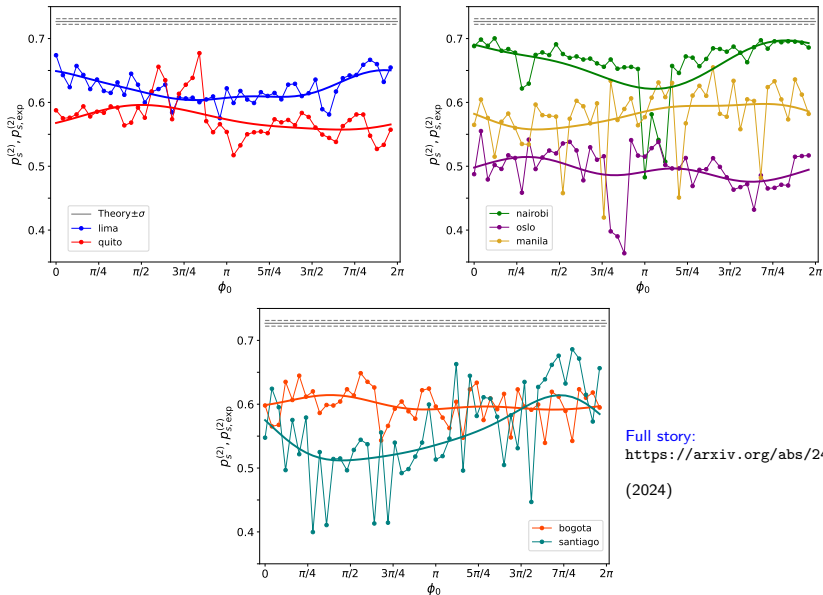


Figure: Success probability  $p_s$  as a function of  $\phi_0$  for  $\theta = \pi/8$  and one step of the protocol.

## Quantum benchmark: Results II - Coherent errors



Full story:  
<https://arxiv.org/abs/2410.07056>  
 (2024)

Figure: Success probability  $p_s$  as a function of  $\phi_0$  for  $\theta = \pi/8$  and one step of the protocol.

## Quantum benchmark: Results III - Statistical metrics

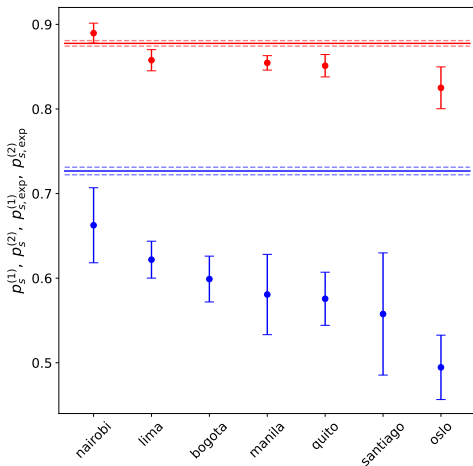
The performance of a quantum circuit in a certain qc is measured with the metrics:

$$F^{(n)} = 1 - \frac{|\bar{p}_{s,\text{exp}}^{(n)} - p_s^{(n)}|}{p_s^{(n)}}, \quad (1)$$

$$S^{(n)} = \frac{\sigma_{s,\text{exp}}^{(n)}}{\sigma_s^{(n)}}, \quad (2)$$

- $F^{(n)}$  estimates how much all measurements are close to the theoretical value of  $p_s$ .
- $S^{(n)}$  measures the “dispersion” in experiment measurements with respect to the pure statistical noise. Captures the total “noise” that affects the qc.

## Quantum benchmark: Results III - Statistical metrics

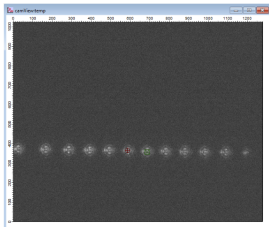


**Figure:** Success probability  $\overline{p}_{s,\text{exp}}^{(n)}$  averaged over the initial angle  $\phi_0$  and its respective variance, with  $n = 1, 2$  steps.

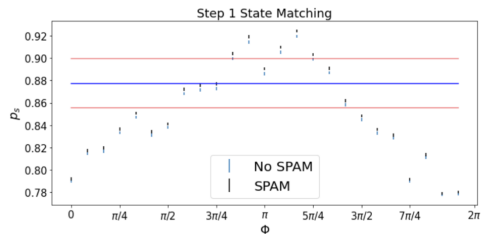
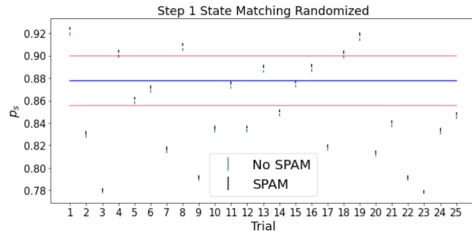
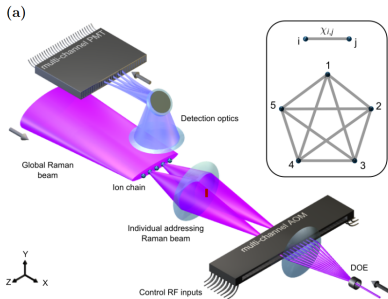
Device	QV	$F^{(1)}$	$S^{(1)}$
nairobi	32	0.986	3.609
lima	8	0.977	3.812
manila	32	0.974	2.599
quito	16	0.970	4.041
oslo	32	0.940	7.544

Device	QV	$F^{(2)}$	$S^{(2)}$
nairobi	32	0.912	9.953
lima	8	0.856	4.893
bogota	32	0.824	6.077
manila	32	0.799	10.650
quito	16	0.792	7.045
santiago	32	0.767	16.204
oslo	32	0.681	8.539

# Quantum benchmark: Results IV - Ion traps (in progress)



← Real ions!



Collabs:

N. Linke, A. Green & E. Mossman





## Outlook

Advantages of our algorithm:

- Highly structured circuit (fast transpilation).
- **No classical overhead** [c.f. Quantum Volume].
- $2^n$  qubits tested [Testing advantage!]
- Opportunity to test any pair of basis states  $|i_0 i_1 \dots i_{2^n-1}\rangle$ ,  
 $|j_0 j_1 \dots j_{2^n-1}\rangle$  [albeit SWAP gates must be introduced...].

Experiments:

- Time variations via the scanning of  $(\phi_0, \theta_0)$  and the invariance of  $f(z)$  over  $\phi_0$ .
- Dependence of the initial states. Random initial states are better?
- Equal Quantum Volume does not imply equal application of  $f(z)$  (!).

¡Gracias por su atención!